

GLOBE

Global Object Exchange

A dynamically fault-tolerant and dynamically scalable
distributed tuplespace for heterogenous, loosely coupled networks

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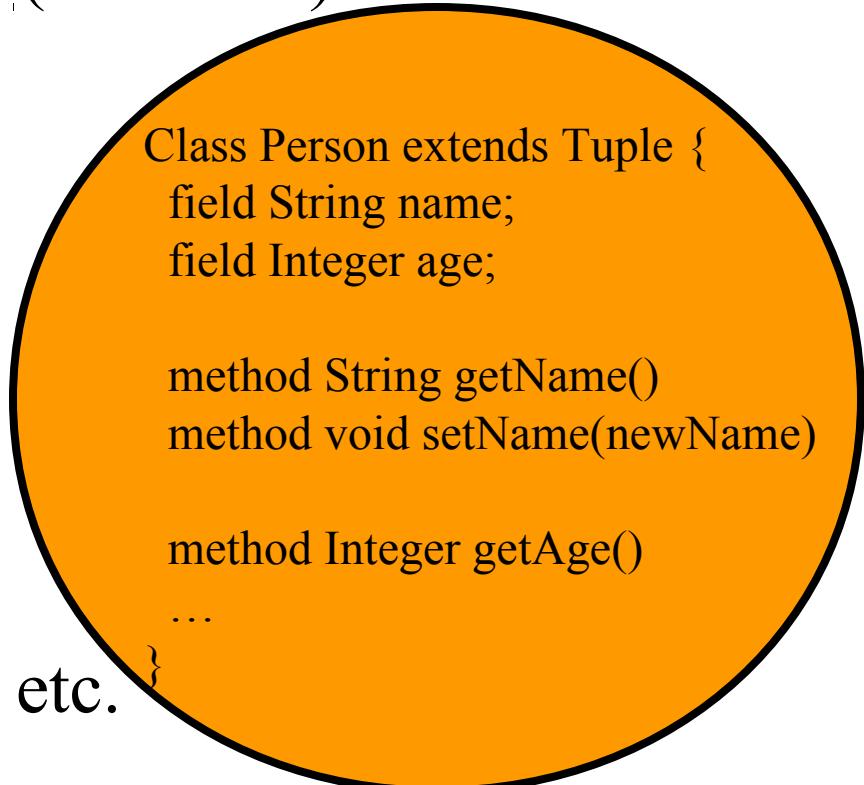
<http://www.diku.dk/students/eglarsen/GLOBE>

Agenda

- The Tuplespace Paradigm
- Tuplespace Semantics
- Achieving Fault-tolerance
- Achieving Scalability
- Measurements of GLOBE
- Related projects
- Conclusion
- Demonstration and Questions

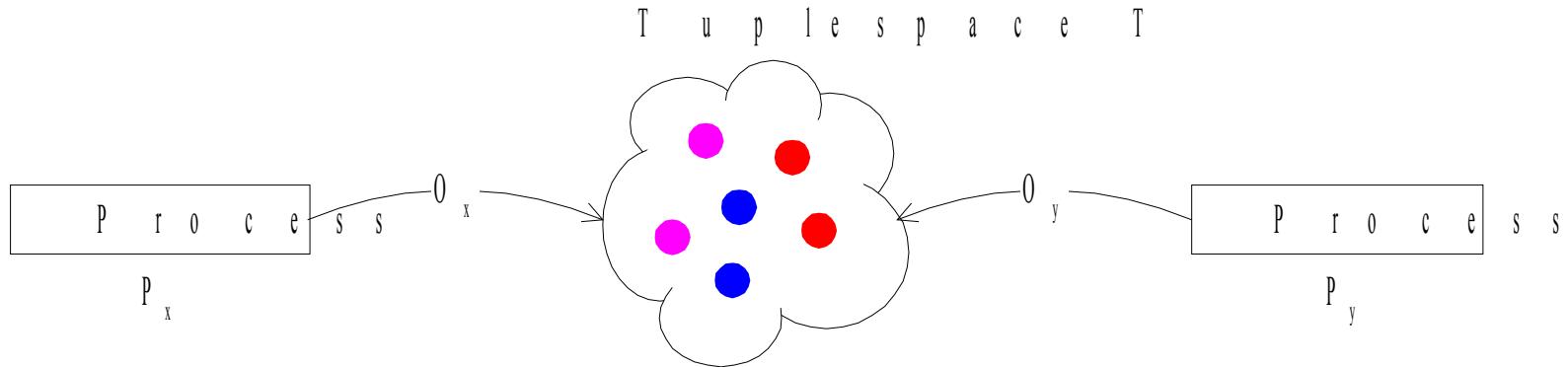
The Tuplespace Paradigm

- A communication paradigm (from '85)
- Communication unit:
Tuple (\approx Object)
- Tuplespace – an intermediate container
- Tuples immutable in the tuplespace
- Implementations:
Linda, JavaSpaces, TSpaces etc.



```
Class Person extends Tuple {  
    field String name;  
    field Integer age;  
  
    method String getName()  
    method void setName(newName)  
  
    method Integer getAge()  
    ...  
}
```

The Tuplespace Paradigm



- Atomic Operations:
 - Insertion (out)
 - Withdrawal (in, inp*)
 - Inspection (rd, rdp*)
 - Additional operations
 - Matching:
 - Templates (anti-tuple)
 - Exact/Wildcard matching:
 - Tuple Type (null tuple)
 - Tuple Field Values (wildcard fields)
- *rdp/inp – predecate

The Tuplespace Paradigm

- Groupware
 - Chat server
 - Shared Blackboard
- High Performance Computing
 - SETI@Home-like calculations
- Intelligent Connectionware
 - Internet Services (internally in Jini LUS)
 - Intelligent Home (TSpaces at IBM)

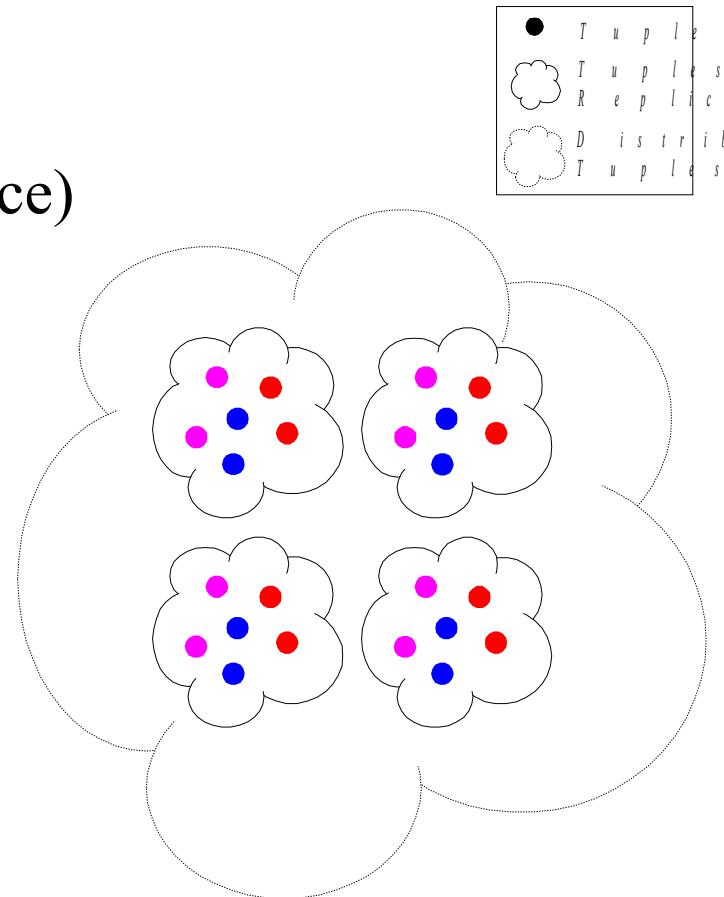
The Tuplespace Paradigm

The Classical Problems with Centralized Systems:

- Availability (Level of Fault-Tolerance)
 - Single Point of Failure
- Scalability
 - Cannot Scale beyond the Single Entity

The Tuplespace Paradigm

- Purpose of GLOBE:
 - Increase Availability (Fault-tolerance)
 - Increase Scalability
 - Dynamic Adjustment
- Distributed Tuplespace Abstraction



Operations Semantics

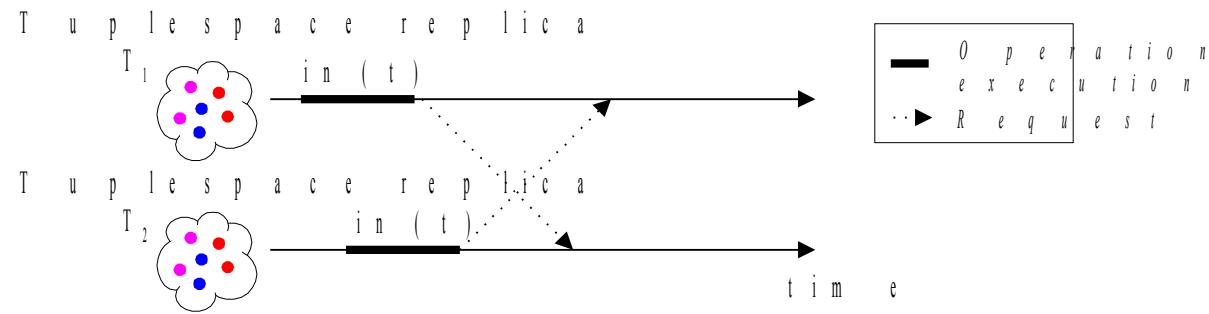
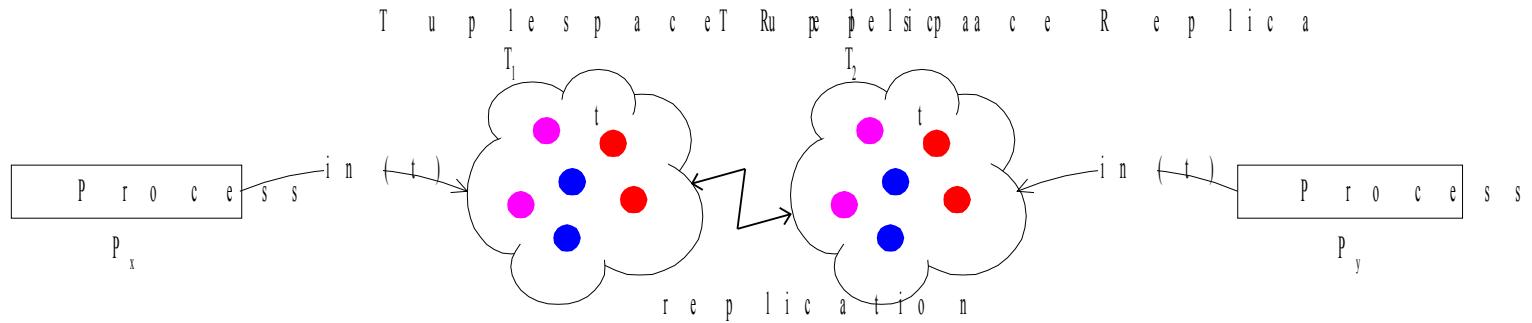
- Tuplespace Semantics vaguely defined:
 - Selection of matching tuple
 - Arbitrary, FIFO, LIFO etc.
 - Selection of process to withdraw tuple
 - Concurrent withdrawals for same tuple
 - Specified as “fair”

In addition: Issue related to distribution:

- Predicate operations semantically unclear
 - Inspect “present” state
 - What is present state in a distributed environment?

Operations Semantics

Concurrent and Distributed Tuple Withdrawal



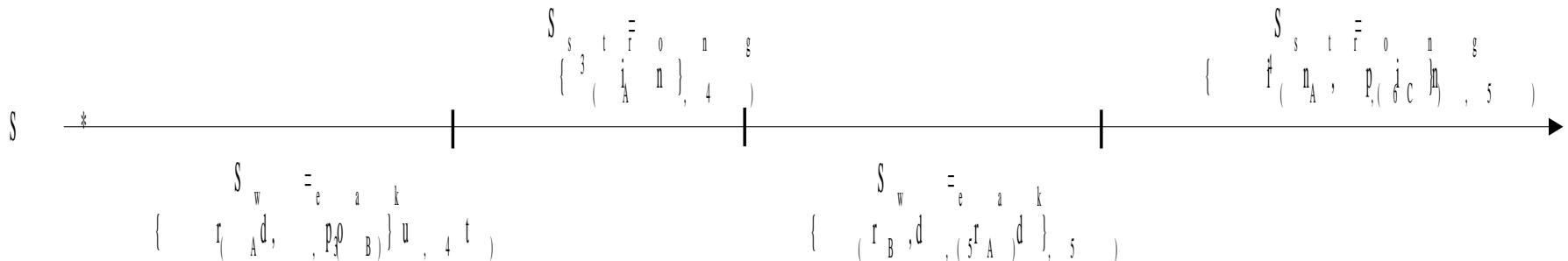
- Withdrawal operations must be performed atomically!
- Global ordering across replicas

Semantics of GLOBE

- Two Categories of Tuplespace Operations:
 - Strong Operations (in, inp)
 - Weak Operations (out, rd, rdp)
- GLOBE adapts *loose inp/rdp* semantics
 - Weak Operations are performed “locally” and any changes (insertions) propagated later (depending on the synchronization tightness).
 - inp/rdp may show “false” results!

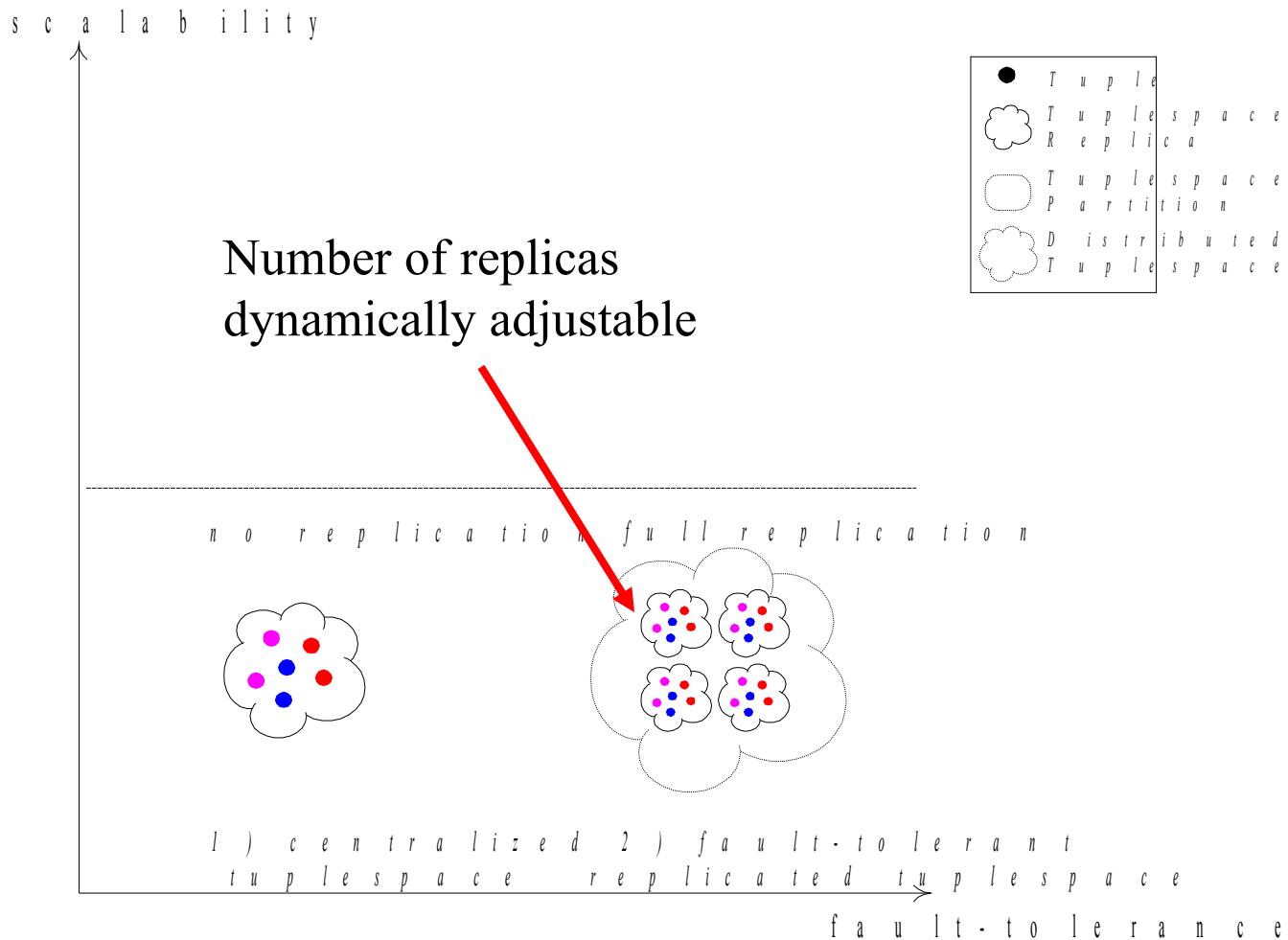
Semantics of GLOBE

- Operation Ordering:
 - Strong Operations are Globally Ordered
 - Weak Operations are Globally Unordered
 - All Operations satisfy Partial Ordering



S^* -- a sequence of tuplespace operations performed on a replica

Achieving Availability (Fault-tolerance) and limited Scalability by Replication



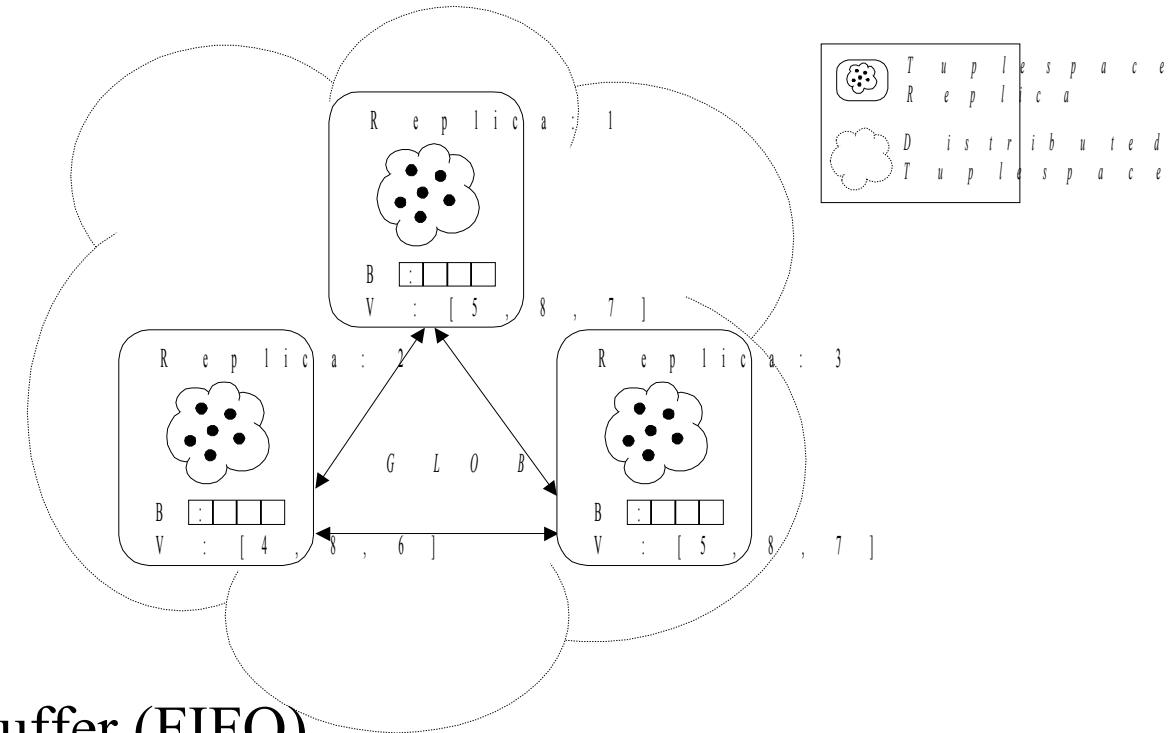
Replica Update Protocol

- Active propagation
 - Operations are persisted before propagation
 - Replicas responsible for propagation
 - Fast propagation
 - Majority voting for all atomic operations (tuplespace operations and configuration operations)
- Problem: Not 100% reliable!
 - Adjustability problem (removal of replica)
 - Inconsistency in case of failure

Replica Update Protocol

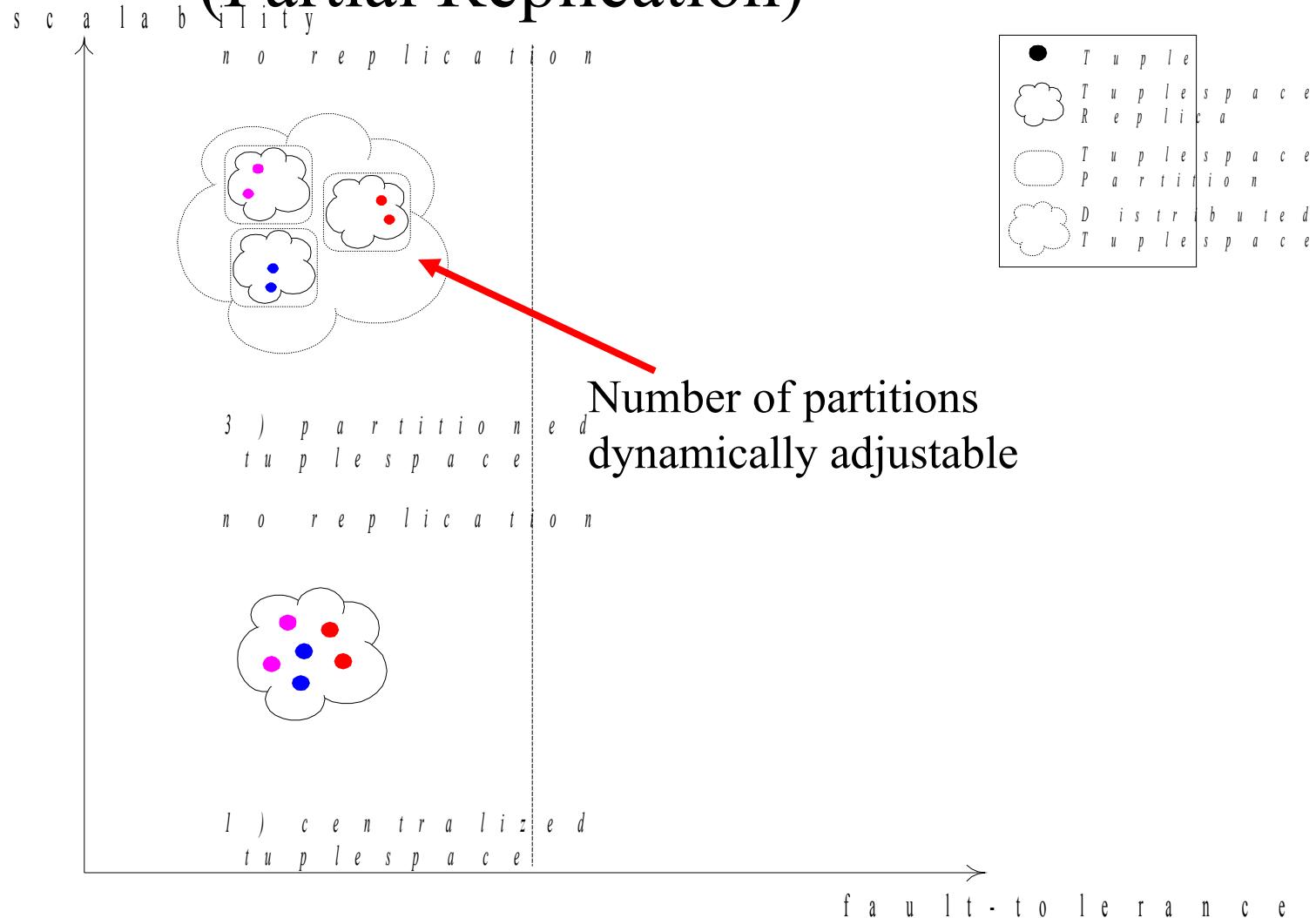
- Anti-entropy
 - Epidemic algorithm (Bayou), pair-wise reconciliation - slow!
 - Replicas responsible for updating themselves (pull-based to avoid duplicates)
 - Synchronization in case of failure
- Hybrid replica update protocol
 - Combines Active Propagation and Anti-Entropy
 - Ensures consistency convergence

Update Propagation



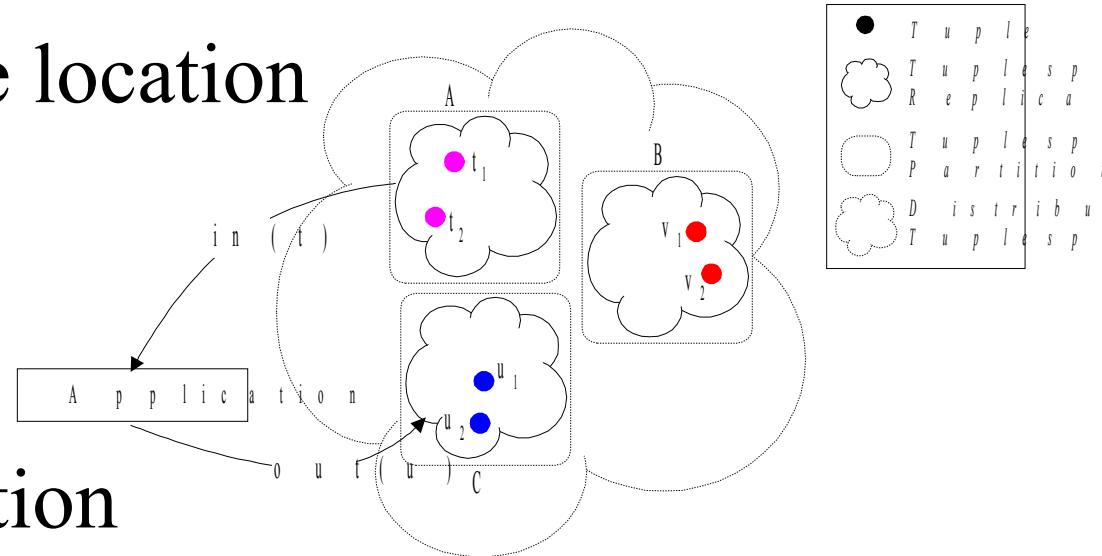
- Operation Buffer (FIFO)
- Operation Vector
- Logical Operation Numbers

Achieving Scalability by Partitioning (Partial Replication)



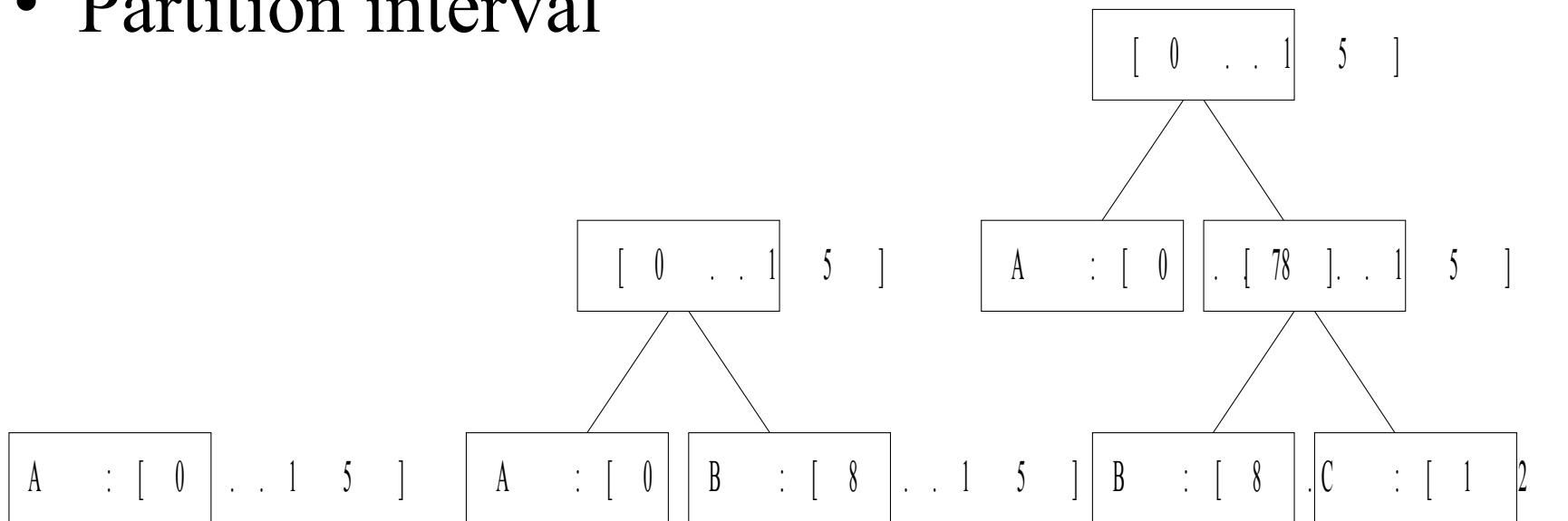
Tuplespace Partitioning

- Load balancing
- Resolution of tuple location
 - Non-deterministic
 - Deterministic
- Hash code
- Operation Redirection
- Problem: Dynamic adjustment
 - Complete rehashing of tuples



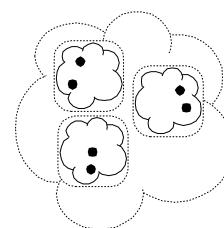
Dynamic Partitioning

- Partitioning by hash code
- Partition interval

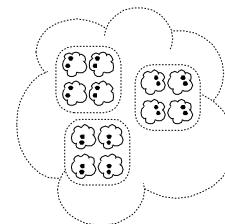


Related Projects

S c a l a b i l i t y



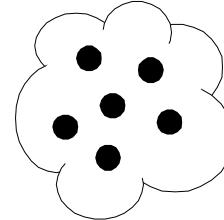
[B j o 9 3]



G L O B E (l o o s e l)
[K r i 9 1] (h a r d w a r e - b a s e d)

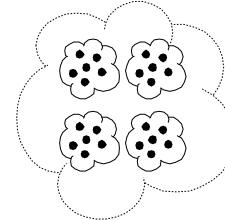
| | | |
|---|-----------------------|-----------------------|
| ● | T u p p l e | s p a c e |
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| ○ | T u p p l e s p a c e | P a r t i t i o n |
| ○ | D i s t r i b u t e d | T u p p l e s p a c e |
| | | e d n e t w o r k |

3) P a r t i t i o n e d f t a u p l l t e - s p a c e c r e a n t t u p p l e s p a c e



O r i g i n a l L i n d
[S u n 9 9 b] J a v a
[W + 9 L 8 b] T S p a c e

I) C e n t r a l i z e d



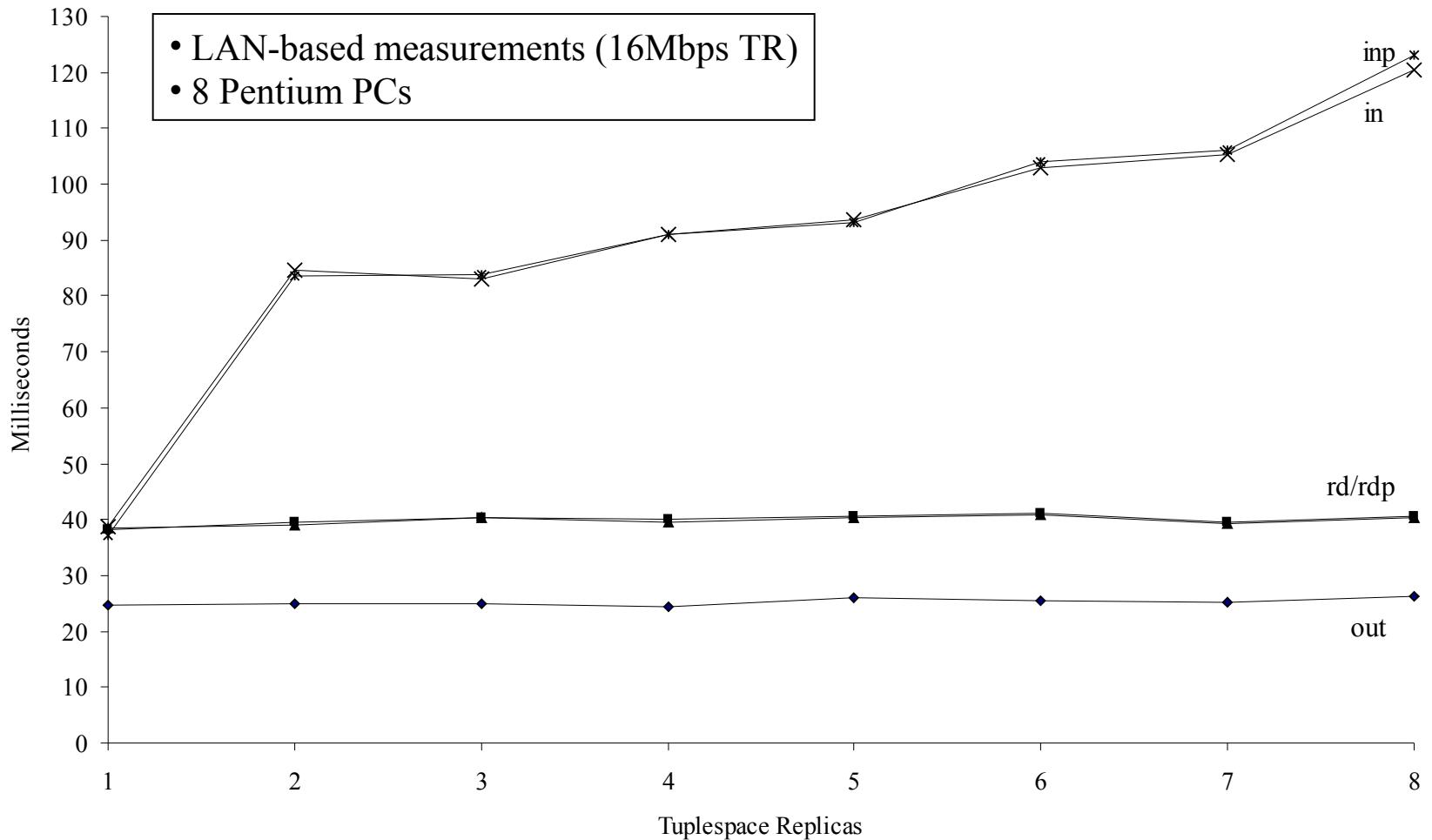
a [C i mG p 8 16] m S e / N t a t i o n s d a
[S B p Sa 8 e 5 s] F T - L i n d a
[sX L 8 9]
[K a m 9 1]
[C K M 9 2] M T S

d 2) u F p a l e u s l p - a t a k e r a n t t u p p l e s p a c e

F a u l t - t o l e r a n c e

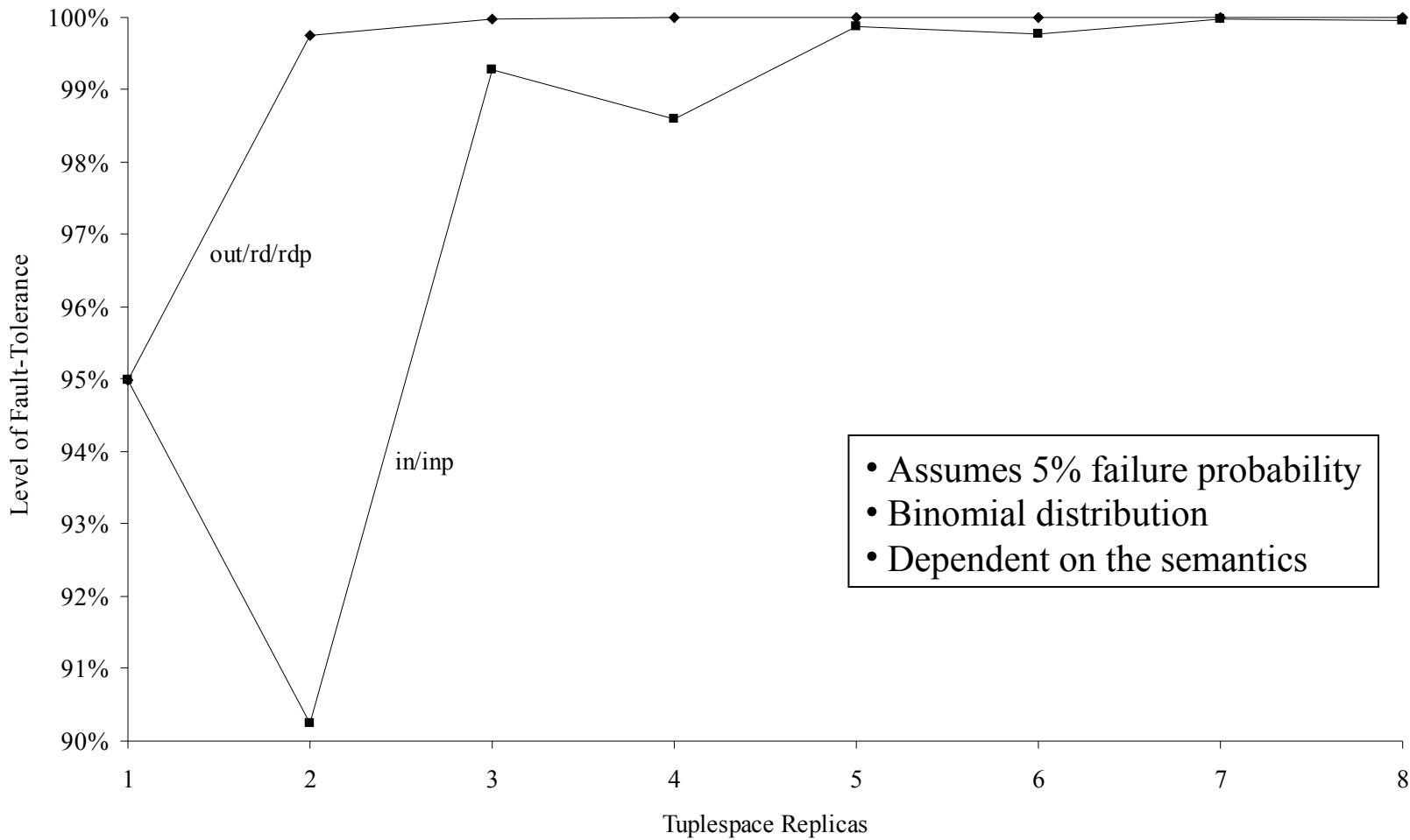
Cost of Fault-Tolerance

(Empirical)



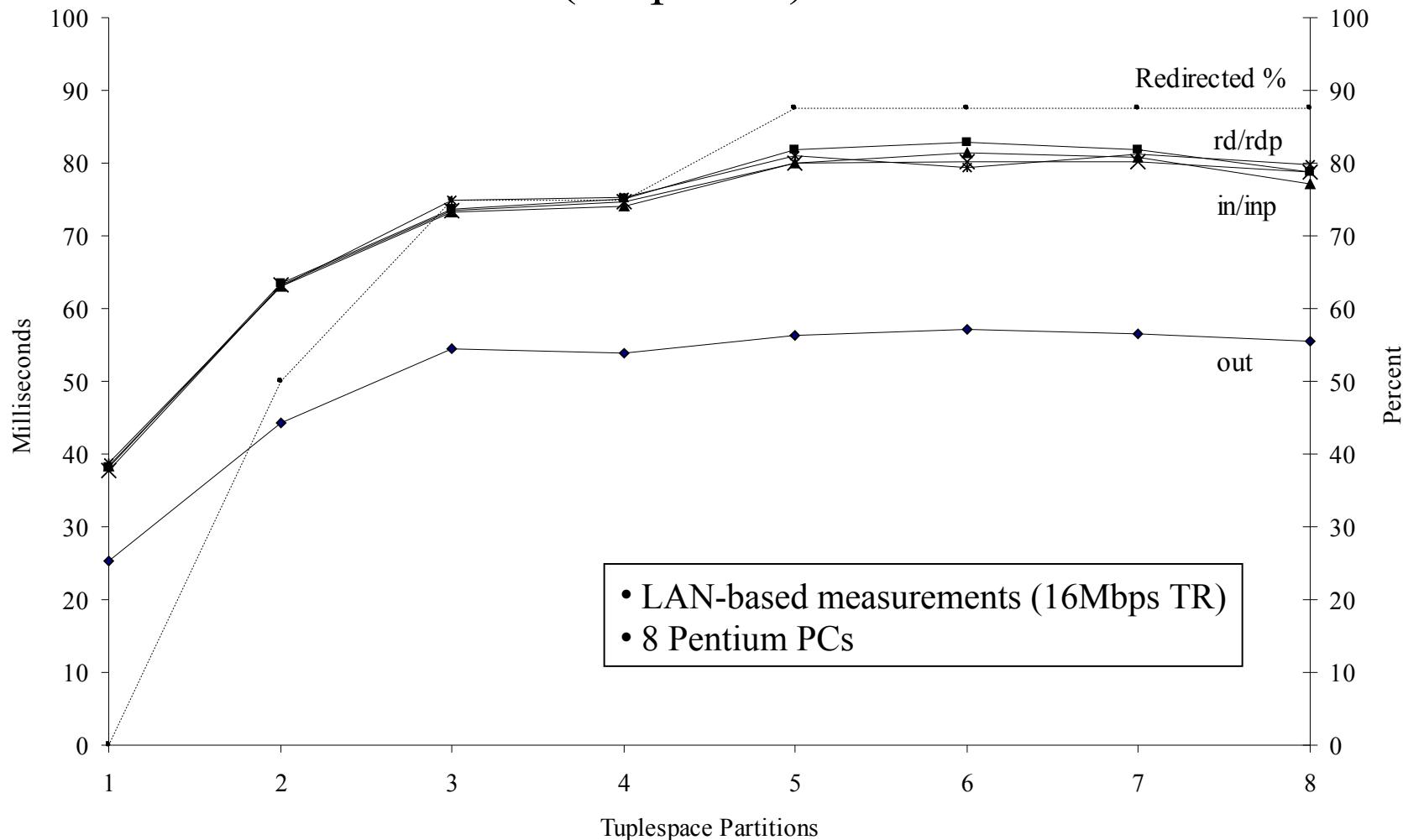
Gains of Fault-Tolerance

(Theoretical)



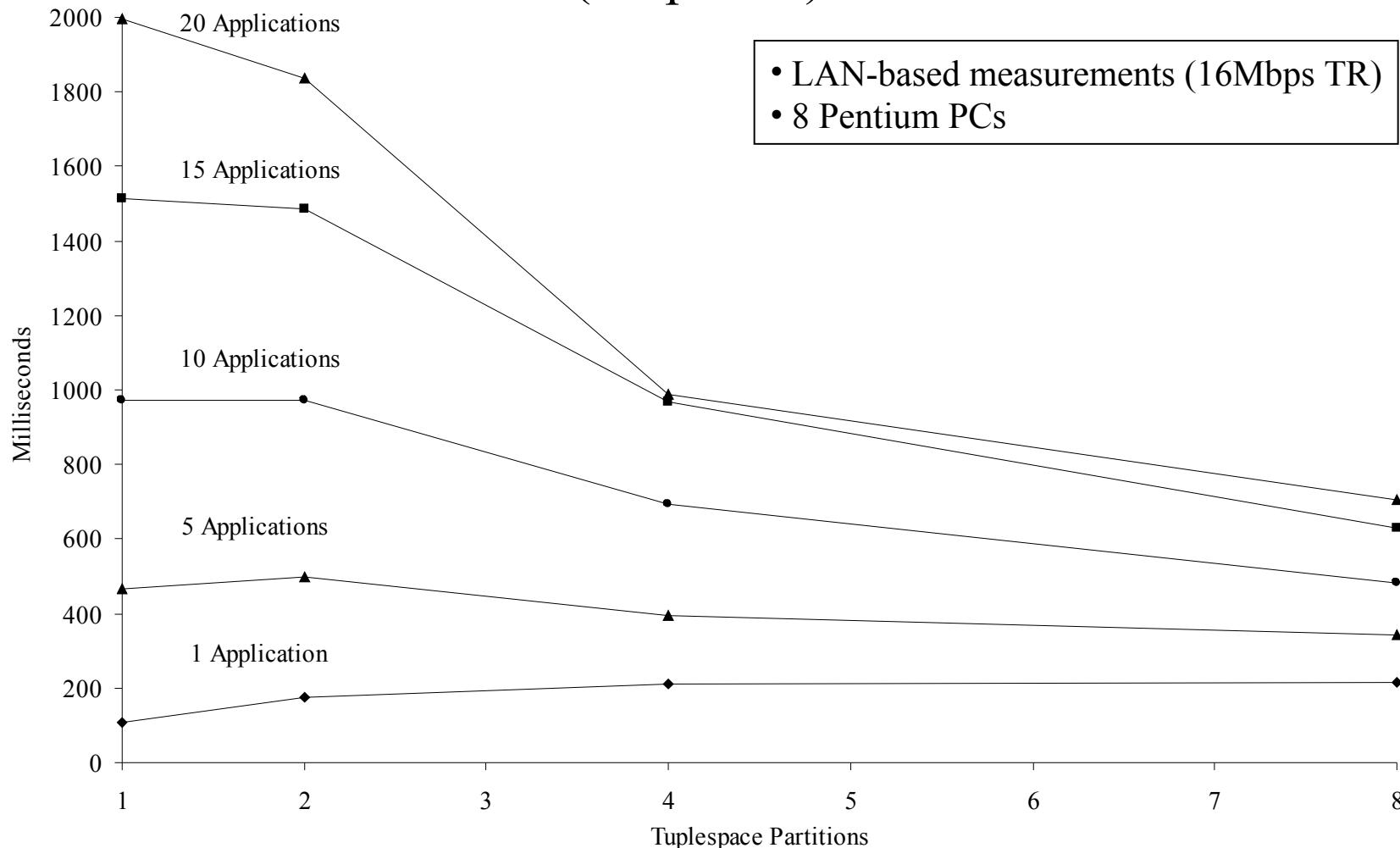
Cost of Scalability

(Empirical)



Gains of Scalability

(Empirical)



Conclusion

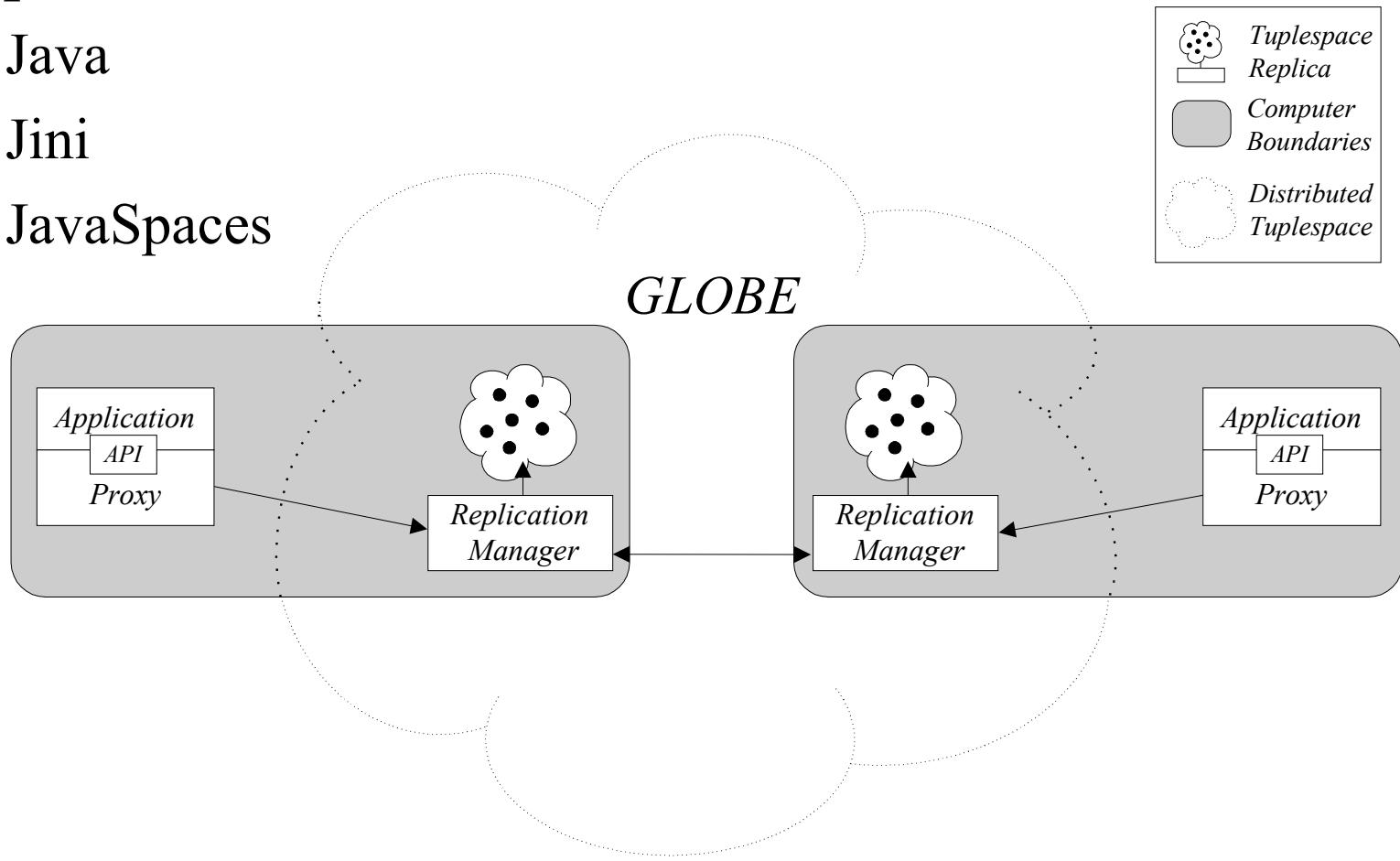
- Distributed Tuplespace Semantics is suitable
- Hybrid Replica Update Protocol is fast and ensures consistency
- Higher level of Fault-Tolerance
- Higher level of Scalability
- Fault-Tolerance and Scalability dynamically Adjustable
- Outperforms a highly loaded centralized Tuplespace

Future Work

- GLOBE enhancements
 - Elimination/reduction of Redirection
 - Load-balancing of Applications
 - Implementation Optimizations
 - Additional Tuplespace Features

Demonstration

- Implementation:
 - Java
 - Jini
 - JavaSpaces



Fault-Tolerance & Scalability

